Name Table

What's in a name? Wed, Sep 9, 1992

The VME local station systems have names that are used for analog channels. They may also have names for binary status bits. For a station with 2K analog channels, the time to execute the current linear search for a 6-character name can be 5 ms. This note describes an imple mentation of a generalized name table used for fast searching of names using a double hashing algorithm.

Initialization

At reset time a table is built that contains all the known names in the system. Each system table (in non-volatile memory) that houses names of anything is scanned, and the names are entered into a hashing table. Let each name length be even and up to 32 characters in length. The name is reduced to a longword and divided by the number of allocated table entries which is the larger prime number of a prime pair. In this way, the quotient is an index into the hashing table. If that entry is examined and found not empty, it is tested for a match against the search name. If there is a match, the name is already in the table. If there is no match, the next entry in sequence is examined, using as a sequencing delta a value obtained by dividing the same dividend by the smaller of the prime pair and adding 1. (This scheme is called double hashing, as it uses two hashing functions; it is described by Knuth in his volume on Sorting and Searching.) This matching procedure is continued in this way until an empty entry is found. Enter the new data into the empty entry. Only when the table is full will there be no empty entry found, although the efficiency of this scheme falls off somewhat before that happens.

Hash table entries

Suppose the information in the entry is 8 bytes as follows:



The del byte is a "deleted" flag. The type byte denotes the name type category, allowing for names of different types to match. The index word is the result data of a name lookup of a given type. (The most obvious example is an analog chan nel number.) The ptr to name is the address of a name field in a system table entry. To allocate room for 8K names, 64K bytes is required for the table. It should be made rather larger than this to reduce the likelihood of collisions.

To process a request that specifies a name ident, the name table is consulted. If the hash code points to a non-empty entry, match against the name field pointed to. If there is no match, continue with the next entry until a match or an an empty entry is found. In the latter case, there is no match on the name, so the name does node# index

Insertions and deletions

To enter a new name into the table, follow the procedure described under Initialization. To delete an entry, the same sequence is followed, but upon finding an match, mark the entry deleted. It cannot me marked empty, since doing so may mask further entries in the chain.

Error return codes: Routine:

0 No errors

Invalid name table (All)
 #chars must be even and in range 2-32 (All)
 All

-3 Name not in table NTLookup, NTDelete

Table overflow
 Duplicate name already in table
 Bus error. Bad ptr in table entry

NTInsert
All

Name changes

When a setting to an analog descriptor is made, a special check is made to determine whether the setting will result in changing the name field for that channel. If so, the current name is deleted and the new name added. This can cause table entries to be used up as entries are deleted, but they can be re-used. At reset time, rebuilding the table removes such "deleted" entries.

Name table header layout

Time to enter 16-char names (0.5 ms)

Total time to initialize name table

NTTIM16

NTTIMT

spare(11)

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